

t-Pebbling Number of Some Multipartite Graphs

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Abstract

Given a configuration of pebbles on the vertices of a graph G , a *pebbling* move consists of taking two pebbles off some vertex v and putting one of them back on a vertex adjacent to v . A graph is called *pebbleable* if for each vertex v there is a sequence of pebbling moves that would place at least one pebble on v . The *pebbling number* of a graph G , is the smallest integer m such that G is pebbleable for every *configuration* of m pebbles on G . The concept of t -pebbling in graphs is the general formulation of pebbling in graphs. In this chapter we characterise the t -pebbling number of a large class of diameter two graphs. We have also established some bounds for the t -pebbling number of graphs.

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Introduction

Pebbling: F R K Chung [2] introduced the concept of pebbling in graphs. A pebbling configuration C of a graph G is a distribution of pebbles on G . A pebbling move consists of removing two pebbles lying on the same vertex v and placing one of these pebbles on some vertex that is adjacent to v . For a given pebbling configuration C , a vertex v is called pebbleable if there is a sequence of pebbling moves such that at least one pebble can be placed on v . The pebbling number of a graph G is the minimum number m of pebbles that ensure that every vertex of G is pebbleable, no matter what initial configuration of m pebbles we start with. Let $f(G, v)$ denotes the pebbling number of a vertex v of G , and $f(G)$ the pebbling number of G . Hence,

$$f(G) = \max_{v \in V(G)} f(G, v)$$

t-pebbling: The concept of t -pebbling in graphs is the general formulation of pebbling in graphs. Let G be a connected graph and let $v \in V(G)$. Let $f_t(G, v)$ denote the smallest integer m such that, whatever be the distribution of m pebbles on the vertices of G , t pebbles can be moved to v .

We define t -pebbling number $f_t(G)$ of a graph G to be

$$f_t(G) = \max_{v \in V(G)} f_t(G, v)$$

Sequential join [3]: Let G_1 and G_2 be graphs such that G_1 and G_2 have disjoint vertex sets V_1 and V_2 and edge sets E_1 and E_2 respectively. Their join $G_1 + G_2$ consists of $G_1 \cup G_2$ and all edges joining V_1 with V_2 . For three or more disjoint graphs G_1, G_2, \dots, G_n , the sequential join $G_1 + G_2 + \dots + G_n$ is the graph $(G_1 + G_2) \cup (G_2 + G_3) \cup \dots \cup (G_{n-1} + G_n)$.

Let

(i) $p(u)$ denotes the number of pebbles at the vertex u .

(ii)
$$p(W) = \sum_{v \in W} p(v)$$

(iii) $p = |C|$ = the total number of pebbles used in C .

The following facts are almost folklore.

1. t -pebbling number of graphs

Fact 1.1.

$$f_t(K_1) = t$$

Fact 1.2.

$$f_t(K_2) = 2t$$

Fact 1.3. For any connected graph G , $f_t(G) \leq f_{t-1}(G) + f(G)$. Equality can hold in the above inequality. For example, if $G = P_n, n \geq 3$, $f_t(G) = t2^n = f_{t-1}(G) + f(G)$ for any $t \geq 2$.

More generally,

Lemma 1.4. Let G be a connected graph of order $n \geq 2$. Let u be any vertex in G . Then $f_t(G, u) \geq 2t + (n - 2)$, if $n \geq 4$. Equality holds if $\deg(u) = n - 1$.

Proof. Let u be the target vertex. Let $v \in N(u)$. Let C be a pebbling configuration of $2t + (n - 3)$ pebbles on the vertices of G such that $p(u) = 0, p(v) = 2t - 1, p(w) = 1$, for all $w \in V(G) - \{u, v\}$. Then t pebbles cannot be moved to u . This implies $f_t(G, u) \geq 2t + (n - 2)$. Now, suppose that $\deg u = n - 1$, and C is a configuration of $2t + n - 2$ pebbles on $V(G)$. Let $p(u) = k$. If $V_1 = V(G) - \{u\}$, then there are $(2t + (n - 2) - k) - (n - 1) = 2t - 1$

extra pebbles on V_1 . Hence $\left\lceil \frac{2t - 1 - k}{2} \right\rceil$ pebbles can be moved to u . Hence after a pebbling process, there will be $k + \left\lceil \frac{2t - 1 - k}{2} \right\rceil = t + \left(k - \left\lfloor \frac{k + 1}{2} \right\rfloor \right) \geq t$ pebbles on u . Therefore, if $\deg u = n - 1$, $f_t(G, u) = 2t + (n - 2)$. Hence the lemma is proved. ■

As a corollary we have the following.

Corollary 1.5. $f_t(K_n) = 2t + (n - 2)$, if $n > 1$.

Hence, unless otherwise specified, hereafter we assume G as a non-complete graph. In particular that $|V(G)| \geq 3$.

Lemma 1.6. Let v be any vertex of degree less than $(n - 1)$, in a connected graph G , of order n . Then $f_t(G, v) \geq 4t$.

Proof. Since $\deg(v)$ less than $(n - 1)$, there exists a vertex u of G such that $u \notin N[v]$. Consider the configuration C with $p = 4t - 1$ such that $p(u) = 4t - 1$, $p(x) = 0$, for every $x \neq u$. Then t pebbles cannot be moved on to v and hence $f_t(G, v) \geq 4t$. ■

Corollary 1.7. For any vertex v of degree less than $n - 1$, we have $f_t(G, v) \geq \max\{4t, 2t + (n - 2)\}$.

A class of graphs for which equality holds in the above inequality is considered in corollary (1.9).

We next consider the t -pebbling number for graphs of the form $G_1 + G_2$.

Theorem 1.8. Let G be a graph with vertex set V which can be partitioned into two nonempty subsets V_1 and V_2 which satisfy the following properties.

- (i) $V_1 = \{u_1, u_2, \dots, u_{r_1}\}$ is an independent set .
- (ii) $V_2 = \{v_1, v_2, \dots, v_{r_2}\}$ where r_1 and r_2 are integers ≥ 2 .
- (iii) $N[u_1] = N[u_2] = \dots = N[u_{r_1}] = V_2$. Then $f_t(G, u_1) = \max\{2t + n - 2, 4t + r_1 - 2\}$.

Proof. We have seen already that $f_t(G, u) \geq 2t + (n - 2)$ for any vertex $u \in V(G)$, for a graph G . Again consider the configuration C , given by $p(u_1) = p(V_2) = 0$, $p(u_2) = 4t - 1$, $p(u_i) = 1$, for all $i, 3 \leq i \leq r_1$. Clearly we cannot move t pebbles to u_1 starting from p . Hence $f_t(G, u_1) \geq \max(2t + n - 2, 4t + r_1 - 2)$. Next, we prove the reverse inequality.

Let $V_3 = V_1 - \{u_1\}$. We prove the theorem by induction on t . When $t = 1$, the theorem reads as $f(G, u_1) = \max\{n, r_1 + 2\} = n$, since $r_1, r_2 \geq 2$ and $r_1 + r_2 = n$. Let C be a configuration of n pebbles on the vertices of G . If $p(u_1) > 0$ or $p(v_i) \geq 2$, for any vertex $v_i \in V_2$, u_1 can be pebbled. So, we may assume $p(u_1) = 0$ and $p(v_i) \leq 1$, for

every $v_i \in V_2 \Rightarrow p(V_3) \geq n - r_2 \geq r_1$. As $|V_3| = r_1 - 1$, there exists at least one $u_i \in V_3$ such that $p(u_i) \geq 2$. If $p(v_i) \geq 1$ for atleast one i , we can move a pebble at u_1 using the path $u_2v_iu_1$. Otherwise, $p(V_2) = 0 \Rightarrow p(V_3) = n$. Since $n(r_1) = r_2 + 1 \geq 3$, atleast 3 excess pebbles are in V_3 . Hence, two pebbles can be moved to V_2 and then one pebble can be moved to u_1 . Hence $f(G, u_1) = n$. Therefore the result is true for $t = 1$. We may assume that $2 \Leftrightarrow n - r_1t \Leftrightarrow 2t \leq r_2$. We may assume $t \geq 2$. We note that $2t + n - 2 \geq 4t + r_1 - 2 \geq 2$.

Case 1(a) : Let $2t \leq r_2$.

In this case, we prove $f_t(G, u_1) = 2t + n - 2$ by induction on t . We have proved the result for $t = 1$. Assume that the result is true for $2, 3, \dots, (t - 1)$. Let p be a distribution of $s = 2t + (n - 2)$ pebbles on G . If either $p(u_1) > 0$ or $p(v_i) \geq 2$, for any vertex $v_i \in V_2, u_1$ can be pebbled at the cost of atleast two pebbles. Then by induction we can move $(t - 1)$ more pebbles to u_1 . Therefore, we may assume $p(u_1) = 0$ and $p(x) \leq 1, \forall x \in V_2$. Let $k = |\{x \mid x \in V_2, p(x) = 1\}| \Rightarrow p(V_3) = s - k$.

Case (1) : $k \geq t$.

In this case, $(s - k) = 2t + [(n - 2) - k] \geq 2t$. We can move one pebble each to the t occupied vertices in V_2 . Then we can move t pebbles to u_1 .

Case (2) : $k < t$.

In this case, if we can move $k + 2(t - k) = 2t - k$ pebbles to V_2 , we are through. For that, we can move the pebbles in such a way that each of the k occupied vertices has an even number of pebbles, the total number being $2t$. From there t pebbles can be moved to u_1 . But $s - k = 2t + (n - 2) - k \geq 4t - k \geq 4t - 2k$. Thus, it is possible to move t pebbles to u_1 starting from p .

Case 2(b) : $2t = r_2 + 1$.

Let C be a configuration of $s = 4t + r_1 - 2$ pebbles on G . We note that $p = 4t + r_1 - 2 = 2t + r_1 + r_2 + 1 - 2 = 2t + n - 1$. Also, as $2(t - 1) < r_2$, by case 1, $f_{t-1}(G, u_1) = 2(t - 1) + n - 2 = 2t - 3$. Consider the following cases,

- (i) $p(u_1) \geq 1$
- (ii) $p(u_i) \geq 2$ for some $u_i \in V_3$ and $p(v_i) \geq 1$ for some $v_i \in V_2$.
- (iii) $p(v_i) \geq 2$ for some $v_i \in V_2$.

In each of the above cases we can move one pebble to u_1 at the cost of atleast 3 pebbles, then we can move $(t - 1)$ pebbles to u_1 . We may therefore assume that none of the preceding conditions holds.

This implies $p(u_1) = 0, p(V_2) \leq r_2$ and $p(V_3) \geq 4t + r_1 - 2 - r_2 \geq 2t + (r_1 - 1)$. If now $p(v_i) = 1$ for some $v_i \in V_2$, we can move one pebble to u_1 at the cost of three pebbles. Then we can move $(t - 1)$ more pebbles to u_1 . So, we may

assume further that $p(V_2) = 0$. Then $p(V_3) = 4t + r_1 - 2$. There are atleast $(4t + r_1 - 2) - (r_1 - 1) = 4t - 1$ extra pebbles on the vertices in V_3 . Hence t pebbles can be moved to u_1 from V_3 through a vertex in V_2 .

Case (c) : We prove in general that, if $t \geq \left\lceil \frac{r_2}{2} \right\rceil$, $f_t(G, u_1) = 4t + r_1 - 2$. We prove the result by induction on t . We have already proved the result when $t = \left\lceil \frac{r_2}{2} \right\rceil$. So, we assume that $f_r(G, u_1) = 4r + r_2 - 2$ if $\left\lceil \frac{r_2}{2} \right\rceil \leq r \leq t - 1$. If $2t = r_2 + 1$ we have proved. When $2t > r_2 + 1$, $2(t - 1) > r_2$. Whatever be the distribution of $4t + r_1 - 2$ pebbles on the vertices of G , one pebble can be moved to u_1 at the cost of atleast 4 pebbles. There are atleast $4(t - 1) + r_1 - 2$ pebbles remaining with $2(t - 1) \geq r_2$. Therefore the result is true by induction on t . The theorem is proved. ■

As mentioned earlier, we consider a class of graphs for which the equality holds in the inequality obtained in corollary (1.7).

Corollary 1.9. Let G be a connected graph of order $n \geq 4$. Let u and v be non adjacent vertices of G such that $\deg u = \deg v = n - 2$, then $f_t(G, v) = \max\{4t, 2t + (n - 2)\}$.

Proof. If $G_1 = \langle u \rangle$, $G_2 = G - \{u, v\}$ and $G_3 = \langle v \rangle$, we see that G is isomorphic to the sequential join $G_1 + G_2 + G_3$. ■

Corollary 1.10. Let K_{m_1, m_2, \dots, m_r} be a complete r -partite graph. Let $n = \sum m_i, \forall i, 1 \leq i \leq r$. Then $f_t(G) = \max\{(2t + n - 2), \max (4t + m_i - 2)\}$, where $1 \leq i \leq r$.

Next we consider the other extreme $G = G_1 + G_2$, where G_1 is complete.

Theorem 1.11. Let $G \cong K_{n_1} + K_{n_2} + K_{n_3}$, then

$$f_t(G) = \begin{cases} 2t + n - 2, & \text{if } n_2 \geq 2t \\ 2t + (n - 2) + (2t - n_2), & \text{if } n_2 \leq 2t \end{cases}$$

Proof. Proof of theorem (1.11) is similar to the proof of theorem (1.8). ■

Question 1.12. Find the t -pebbling number of the sequential join of r complete graphs $K_{n_1}, K_{n_2}, \dots, K_{n_r}$, for $r \geq 4$.

2. Bound for the t -pebbling number of graphs

Now, we establish some bounds for the t -pebbling number of graphs.

Theorem 2.1. If G is a connected graph of order n . Let $d =$ diameter of G . Then $f_t(G) \geq \max\{t2^d, 2t + n - 2\}$.

Proof. Let u_0, u_1, \dots, u_d be a diametral path in G . Consider the pebbling configuration C , given by $p(u_d) = (t2^d) - 1, p(v) = 0$, for every $v \in V(G) - \{u_d\}$. Then it is not possible to move t pebbles to $v_0 \Rightarrow ft(G) \geq t2^d$. Again, by lemma 5.3, $f_t(G) \geq 2t + n - 2. \Rightarrow ft(G) \geq \max\{t2^d, 2t + n - 2\}$. The theorem is proved. ■

Chan and Godbole [1] obtained an upper bound for the pebbling number of diameter d graph as $f(G) \leq (n - d)(2^d - 1) + 1$. Using exactly the same argument, we establish an upper bound for the t pebbling number of diameter d graphs.

Theorem 2.2. Let G be a connected graph of order n . Let $d = \text{diam}(G)$. Then $f_t(G) \leq (t - 1)2^d + (n - d)(2^d - 1) + 1$.

Proof. Proof is by induction on t . When $t = 1$, the inequality reads as $f(G) \leq (n - d)(2^d - 1) + 1$, which is true (from [1]). When $t \geq 2$, let $u \in V(G)$. Let e be the eccentricity of u . Let C be a configuration of $(n - e)((t2^e) - 1) + 1$ pebbles on $V(G)$. Let $u_0 = u, u_1, u_2, \dots, u_e$, be an eccentric path with $d(u_0, u_e) = e$. [An eccentric path of a vertex v in G is a path of length $e(v)$ starting from v in G and ending at u where $d(u, v) = e(v)$].

Consider the sets $S_0 = \{u, u_1, u_2, \dots, u_e\}, S_i = \{x_i\}, 1 \leq i \leq n - e - 1$ where the remaining vertices in $G = V(G) - S_0 = \{x_1, x_2, \dots, x_{n-e-1}\}$. Thus, in p , at least one set must receive $(t2^e)$ pebbles and t pebbles can be moved to u . Thus $f_t(G, u) \leq (n - e)((t2^e) - 1) + 1$.

We note that $(n - e - 1)(t2^{e+1} - 1)(n - e)((t2^e) - 1) = (n - e)(t2^e)(t2^e - 1) \geq 1$, if $e < d \leq n - 1$. Thus, $(n - e)((t2^e) - 1) + 1 \leq (n - d)((t2^d) - 1) + 1$. ■

Remark 2.3. We note that equality is attained in the inequality for paths.

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