Comparative Analysis of Different Topologies based on Network-on-Chip Architectures

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Abstract

As the integration density and complexity of the system-on-Chip (SOC) increases, the conventional interconnects are not suitable to fulfil the demands. The application of traditional network technologies in the form of Network-on-Chip is a possible solution. NoC design space has numerous variables. As an improved topology is selected complexities decrease and power-efficiency increases. In this paper, the main research field in Network-on-chip design focussing on optimized topology design is analyzed. The simulation is modelled using a conventional network simulator tool ns-2, in which by selecting proposed Topology 35.7 % reduction in traversing the longest path is observed.

Keywords: NoC, SoC, Routing, Mesh, ns-2, TCL

INTRODUCTION

Recent technological development in the field of integrated circuits has enabled designers to accommodate billions of transistors. The level of integration has enhanced computational power enormously. The exponential decrease in the feature size has enabled integration of heterogeneous IP cores on a single chip leading to a new era of integration circuits known as System-on-Chip. However, as the number of components and their performance continue to increase, the design of power, area and performance efficient communication infrastructure is gaining equal importance. The traditional methods of connecting these heterogeneous IP Cores are not meeting the demands of these very complex structures. Furthermore, with technology scaling, traditional global interconnects cause problems like synchronization errors, unpredictable delays and high power consumption. [1] Traditional bus and crossbar based methods

to communication become very inefficient, resulting in massive numbers of wires, failed timing closure, increased heat and power consumption, and routing congestion leading to increased die area. The Network–on-Chip approach promises the alternative to traditional bus-based and point-to-point communication structures. The networking methods have been dealing with same kind of problems on traditional computer networks. It indicates that NoC designers can borrow the concept of conventional computer networking with necessary customization to suit demands of SoCs.

The SoCs consists of heterogeneous IP-Cores such as Video processors, Image processors, memory blocks etc. Each of these cores is connected to NoC through a network interface or network adapter Module. The NoCs contain a network of routers responsible for end to end delivery of the packets from IP-cores. The communication demands of these IP-cores vary depending on the application Running on it. The network interface provides seamless integration of theses IP-Cores and network. Locating the interconnect logic closest to each IP block results in fewer gates, fewer and shorter wires, and a more compact chip floor plan. Having the option to configure each connection's width, and each transaction's dynamic priority assures meeting latency and bandwidth requirements. The routers are connected to each other through links. The origin of NoC has also been viewed as a paradigm shift from computation centric to communication-centric design as well as the implementation of scalable communication structures. The modular architecture of NoC makes chip structure highly scalable and well controlled electric parameters of the modular block improve reliability.

As the network communication latency depends on the characteristics of the target application, computational elements and network characteristics (e.g. network bandwidth and buffer size [2]. First of all the target applications and their associated traffic patterns and bandwidth requirements for each node in the network is determined. This application partitioning and knowledge of overall system architecture significantly impact the network traffic and helps determine the optimal network topology. Optimal network topology creates immense impact of design cost, power and performance and helps designers to choose effective and efficient routing algorithms and flow control scheme to manage incoming traffic.

The design space of a NoC is very large, and includes topology choice (mesh, torus, star, etc.), circuit switched or packet switched, and other parameters (link widths, frequency, etc.). Because the traffic patterns of most SoCs can be known, a custom generated network topology and physical placement of components yields better performance and power than a regular-pattern network [4]. A NoC's buffers and links can consume near 75% of the total NoC power [5], thus there is significant benefit to optimizing buffer size, link length and bandwidth of a NoC design.

Generally speaking, determining the optimal topology to implement any given application does not have a known theoretical solution. Although the synthesis of customized architectures is desirable for improved performance, power consumption and reduced area, altering the regular grid-like structure brings into the picture significant implementation issues, such as floor planning, uneven wire lengths (hence, poorly controlled electrical parameters), etc. Consequently, ways to determine efficient topologies that trade-off high-level performance issues against detailed implementation constraints at micro- or nano-scale level need to be developed.

BACKGROUND

Network-on-Chip (NoC) is an emerging paradigm using packet switched networks for communications within large VLSI system-on-Chip. NoCs are poised to provide enhanced performance, scalability, modularity, and design productivity as compared with previous communication architectures such as busses and dedicated signal wires. With the emergence of large number of cores in general purpose and system-on-chip (SoC), NoCs are likely to be prevailing on-chip interconnect fabric. [6]

The early work and basic principles of NoC paradigm were outlined in various seminal articles, for example [7-17] and few text books [18-20]. However, the aforementioned sources do not present many implementation examples or conclusions.

Networking concepts from the domains of telecommunication and parallel computer do not apply directly on chip. From a networking perspective, they require adaptation because of the unique nature of VLSI constraints and cost e.g. area and power minimization are essential; buffer space in on-chip switches are limited, latency is very important, etc. At the same time, there are new degrees of freedom available to the network designer, such as the ability to modify the placement of network endpoints. From the view point of VLSI designer, many well understood problems in the real aim of chip development methodology get a new slant when they are formulated for a NoC based system, a new trade-offs need to be comprehended. Therefore, the field offer opportunities for noble solutions in network engineering as well as system architecture, circuit technology, and design automation. [6]

Current complex on-chip systems are also modular, but most often the modules are interconnected by an on-chip bus. The bus is a communication solution inherited from the design of large board- or rack-systems in the 1990's. It has been adapted to the SoC specifics and currently several widely adopted on-chip bus specifications are available [31-34].

While the bus facilitates modularity by defining a standard interface, it has major disadvantages. Firstly, a bus does not structure the global wires and does not keep them short. Bus wires may span the entire chip area and to meet constraints like area and speed the bus layout has to be customized [35]. Long wires also make buses inefficient from an energy point of view [36]. Secondly, a bus offers poor scalability. Increasing the number of modules on-chip only increases the communication demands, but the bus bandwidth stays

the same. Therefore, as the systems grow in size with the technology, the bus will become a system bottleneck because of its limited bandwidth.

Recently, network-on-chip (NoC) architectures are emerging as a candidate for the highly scalable, reliable, and modular on-chip communication infrastructure platform [11]. The NoC architecture uses layered protocols and packet-switched networks which consist of on-chip routers, links, and network interfaces on a predefined topology. There have been many architectural and theoretical studies on NoCs such as design methodology [10], [11], topology exploration [21], Quality-of-Service (QoS) guarantee [22], resource management by software [23], and test and verifications [24].

In large-scale SoCs, the power consumption on the communication infrastructure should be minimized for reliable, feasible, and cost-efficient implementations. However, little research has reported on energy- and powerefficient NoCs at a circuit or implementation level, since most of previous works have taken a top-down approach and they did not touch the issues on a physical level, still staying in a high-level analysis. Although a few of them were implemented and verified on the silicon [25], [26],they were only focusing on performance and scalability issues rather than the power-efficiency, which is one of the most crucial issues for the practical application to SoC design.

Network-on-Chip (NoC) architectures employing packet-based communication are being increasingly adopted in System-on-Chip (SoC) designs. In addition to providing high performance, the fault-tolerance and reliability of these networks is becoming a critical issue due to several artifacts of deep sub-micron technologies. Consequently, it is important for a designer to have access to fast methods for evaluating the performance, reliability, and energy-efficiency of an on-chip network. [27]

While on-chip networks have been proposed and studied in the academic literature, to date there have been very few implementations of routed on-chip networks. Dally and Towles [10] proposed a 2D torus network as a replacement for global interconnect. They claim that on-chip network modularity would shorten the design time and reduce the wire routing complexity. On-Chip routed networks have also been proposed for use in SoCs such as in CLICHÉ [12], in which a 2D mesh network is proposed to interconnect a heterogeneous array of IP blocks.

A performance analysis also shows that dynamic resource allocation leads to the lowest network latencies, while static allocation may be used to meet QoS goals. Combining the power and performance figures then allows an energy-latency product to be calculated to judge the efficiency of each of the network [28].

In his work, Nikolay K. Kavaldjiev, used run-time reconfigurable NoC for streaming DSP applications taking the advantage of a global communication architecture that avoids limitation by structuring and shortening the global wires. He also proposed architecture of a virtual channel router, which in contrast to conventional architectures is able to provide predictable

communication services and has a lower implementation area and cost than conventional architectures. Dynamic reconfiguration is essential to support the dynamically changing demands of the application domain: the system operates in a constantly changing environment. The user demands change (e.g., starting/terminating applications), the environmental conditions change (e.g., available networks, wireless channel conditions) and the available power budget also changes (decreasing battery budget or connected to the mains). The set of running applications and tasks in the system adapts dynamically to these changes. The run-time reconfiguration modifies the system communication demands. For example, a new data stream may be needed or some of the old streams may be redirected or replaced. The NoC must be able to handle such dynamically changing traffic conditions. Run-time changes in part of the traffic must be possible without disturbing the rest of the traffic. The network reconfiguration time must be short enough to enable adequate system reaction time and reconfiguration must be transparent to the user. [30] The major goal of communication-centric design and NoC paradigm is to achieve greater design productivity and performance by handling the increasing parallelism, manufacturing complexity, wiring problems, and reliability. The three critical challenges for NoC according to Owens et al. are: power, latency, and CAD compatibility [17]. The key research areas in Network-on-Chip design can be summarized as [29]: Communication infrastructure: topology and link optimization, buffer sizing, floorplanning, clock domains, power Communication paradigm: routing, switching, flow control, quality of service, network interfaces Benchmarking and traffic characterization for design and runtime optimization Application mapping: task mapping/scheduling and IP component mapping.

METHODOLOGY

Network-on-Chip is a new paradigm for interconnecting today's heterogeneous IP cores based System-on-Chips (SoCs). In SoC's IP Cores are connected to network of routers using network interfaces and network is used for packet switched on-chip communication. Conventional computer design tools i.e. Network Simulator-2 utility are used for network design and simulation. It provides a versatile practice and visualization environment for the design, configuration, and troubleshooting of network environments. The work done by us uses same tool to compare two topologies. The 2-D mesh is currently the most popular regular topology used for on-chip networks in tile-based architectures, because it perfectly matches the 2-D silicon surface and is easy to implement. However, a number of limitations have been proved in the open literature, especially for long distance traffic. In this type of topology, every node has a dedicated point to point link to every other node in the network. This means each link carries traffic only between the two nodes it connects. If N is total no of nodes in network. Number of links to connect these nodes in mesh = N (N-1)/2 Each node should have (N-1) I/O ports as it require connection to every another node. The advantages are:

- No traffic problem as there are dedicated links. Robust as failure of one link does not affect the entire system.
- Security as data travels along a dedicated line.
- Points to point links make fault identification easy.

The disadvantages are:

The hardware is expansive as there is dedicated link for any two nodes and each device should have (N-1) I/O ports.

- There is mesh of wiring which can be difficult to manage.
- Installation is complex as each node is connected to every node.

As earlier studies have shown that maximum power is consumed by links and interconnect infrastructure. Reducing interconnects and links will result in lower power consumption but can also affect the performance and reliability negatively. The topology suggested by us reduces the number of links thus resulting into lower power consumption keeping same level of reliability and performance levels.

SIMULATION

Network Simulator Ns-2 The simulator, ns-2, has facilities to describe network topology, network protocols, routing algorithms and communication traffic generation. It provides basic TCP and UDP as the network transmission protocols, four routing strategies (Static, Session, Dynamic and Manual) and many mechanisms for modelling traffic generation. It is possible to generate a traffic at random, by burst or with bias towards destinations. Additionally, the simulator has the possibility of incorporating protocols, routing algorithms and traffic generation defined by the user.

The simulator is written in C++ and uses OTcl (Object Tool Command Language) for building command and configuration interfaces. The source code of ns-2 is also available[5]. Ns-2 provides well documented trace format for interpreting simulation results. A graphical animator tool, nam (Network AniMator), is also built into ns-2 for user's friendly visualization of the flow of messages and the whole system simulated. In this paper, a generic NOC architecture would be modelled and simulated in ns-2 with only built-in options. Tcl is used for specifying the NOC simulation model and running the simulation.



Figure-1



As shown in figure 1 and figure 2, the number of links in the mesh topology is 24 while in proposed topology the number of links are 20. The number of hops a packets traverses in the longest path is 5 in Figure 1 while 4 in Figure 2. The comparative result in the form of end to end delay and throughput are shown below.





delay of scene-1



Throughput-2

CONCLUSION

The results achieved in terms of time and reduction in number of links displayed here is encouraging and motivates us to take the work further. As discussed earlier the NoC technology can borrow the tools and techniques from conventional computer network technology with required customization. In our future work, we intend to test same on a standard NoC benchmark. The other design parameters on NoC will also be explored.

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