

Gateway Relocation and Admission Control in Mobile WiMAX Networks

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Abstract

The wimax is a standard which define two tired mobility management service. This forum defined the service which leads to minimize the packet loss and handover delay. However, it leads to another problem: When to perform ASN GW (Access Service Network Gateway) relocation? The standards only define the ASN GW relocation procedures without specifying when the ASN GW relocation should be performed. It is left for vendors and operators to develop their own proprietary solutions. In this project, we propose an algorithm, which incorporates traditional Admission Control (AC) and Wiener Process (WP)-based prediction algorithms to determine when to carry out ASN GW relocation.

Keyword: Mobility management, resource management, admission control, WiMAX networks, statistics and stochastic process, and wireless networks.

1. Introduction

The Access Service Network (ASN) provides wireless radio access for WiMAX subscribers. It consists of one ASN Gateway (ASN GW) and many basestations (BSs). Each ASN is connected to Connectivity Service Network (CSN), which provides IP connectivity services. The WiMAX Forum has defined a two-tiered mobility management: ASN Anchored Mobility and CSN Anchored Mobility:

ASN Anchored Mobility refers to the procedures associated with the MS's movement between BSs, which may belong to the same or different ASNGWs. In ASN Anchored Mobility, the context of the designated MS is transferred from the

previous BS to the new BS. Without performing CSN Anchored Mobility, ASN Anchored Mobility can minimize Handover delay and packet loss.

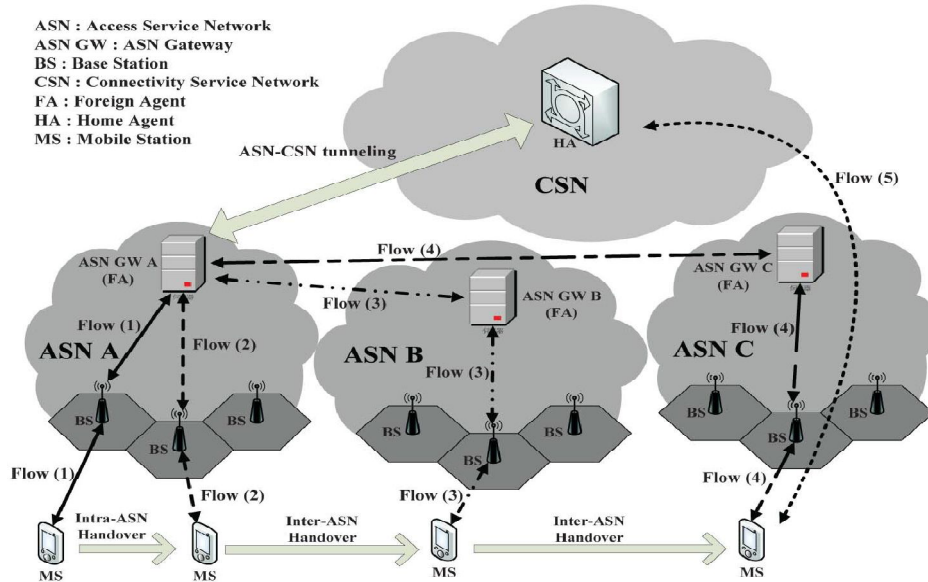


Fig. 1: ASN Anchored Mobility and CSN Anchored Mobility in WiMAX networks.

CSN Anchored Mobility refers to the process of changing the traffic anchor point and is independent of the MS's link layer handover. It is also called ASN GW relocation.

Although the two-tiered mobility management defined in WiMAX potentially can minimize handover delay and packet loss, it leads to another problem: When to perform ASN GW relocation?

In this paper, we propose Gateway Relocation AC (GRAC), which combines ASN GW relocation and AC algorithm to maximize system capacity. In GRAC, the AC algorithm cooperates with the ASN GW relocation. When a new MS arrives and there is no resource for the newly arrived MS, the proposed GRAC will request an Anchored MS to perform ASN GW relocation if there are Anchored MSs in the system.

We also propose a prediction algorithm based on Wiener Process to request Anchored MSs to perform ASN GW relocation early. Thus, handover MSs are not dropped when the system load is full. In addition, handover MSs do not need to wait for the completion of ASN GW relocation so handover latency can be reduced.

2. Related work

Admission Control (AC) is one of the resource management techniques to limit maximum amount of traffic in the network to guarantee service quality for subscribers. In wireless and mobile networks, the AC algorithms are much more complicated due to

the movement of MSs. An MS served in current network may move to another network. The connection of the MS may be dropped if the required resources in the target network cannot be supported. It is generally agreed that keeping an ongoing connection unbroken is more important than admitting a new MS. Here, we discuss two commonly used priority-based AC algorithms: cutoff priority algorithm and new call bounding algorithm.

2.1. Gateway Relocation Admission Control (GRAC)

Our goal is to design a stand-alone algorithm such that each ASN GW can determine when to request Anchored MSs to perform ASN GW relocation. The proposed algorithm does not need to exchange information between neighboring ASN GWs. It also does not require centralized coordination and any assistance from extra servers. In addition, the proposed algorithm does not need to predict the movement of the mobile stations. It combines AC algorithm with a prediction technique to determine when is necessary to perform ASN GW relocation. Thus, it is called Gateway Relocation AC (GRAC).

2.1.1 New Call Bounding AC with ASN GW Relocation

Algorithm for the two-tier mobility management in WiMAX. The proposed GRAC with the new call bounding algorithm is presented in Algorithm given below. In given Algorithm, we limit the number of Serving MSs and Anchored MSs in one ASN GW.

Table 2.1: List of parameters.

C	Maximum number of MSs in one ASN GW
T_{nch}	Threshold for blocking a new MS
T_{wnr}	Threshold for carrying out WP-based prediction
$W(t)$	Number of MSs in one ASN GW at time t
$N_S(t)$	Number of serving MSs in one ASN GW at time t
$N_A(t)$	Number of anchored MSs in one ASN GW at time t
$N_H(t)$	Number of handover MSs in one ASN GW at time t
α	Standard normal random variable
Δt	Prediction time interval
τ	Sampling time interval
k	Number of latest samples
λ_n	Arrival rate of new MSs
λ_h	Arrival rate of handover MSs
$1/\mu_c$	Average connection holding time for new MSs
$1/\mu_n$	Average network residence time for new MSs and handover MSs
p_{nb}^u	Blocking probability of new MSs in upper-bound analysis
p_{nb}^l	Blocking probability of new MSs in lower-bound analysis
p_{hd}^u	Dropping probability of handover MSs in upper-bound analysis
p_{hd}^l	Dropping probability of handover MSs in lower-bound analysis
Θ_u	Average serving rate in upper-bound analysis
Θ_l	Average serving rate in lower-bound analysis
Λ_u	Average signaling overhead in upper-bound analysis
Λ_l	Average signaling overhead in lower-bound analysis

Algorithm 1. New call bounding AC with ASN GW relocation

Require: A new or handover MS is requesting to connect with the ASN GW at time t .

1. if a new MS arrives then
2. if $N_s(t) + N_a(t) < \min(T_{ncb}, C - N_h(t))$ then
3. $N_s(t) = N_s(t) + 1$ /* The new MS is accepted. */
4. else if $N_s(t) + N_a(t) = \min(T_{ncb}, C - N_h(t))$ then
5. if $N_a(t) > 0$ then
6. $N_a(t) = N_a(t) - 1$ /* Requesting one of the Anchored MSs to perform ASN GW relocation. */
7. $N_s(t) = N_s(t) + 1$ /* The new MS is accepted. */
8. else
9. The new MS is blocked.
10. end if
11. end if
12. else if a handover MS arrives then
13. if $W(t) < C$ then
14. $N_h(t) = N_h(t) + 1$ /* The handover MS is accepted. */
15. else
16. The handover MS is dropped.
17. end if
18. end if

To adapt the new call bounding algorithm into WiMAX networks, the algorithm is modified as:

If $N_s(t) + N_a(t) < \min(T'_{ncb}, C' - N_h(t))$ and a new MS arrives; the new MS is accepted.

Because one Anchored MS is relocated, the new MS can be accepted. Otherwise, the new MS is blocked. Further-more, if a handover MS arrives at time t , it is always accepted unless $W(t) = C'$.

2.2.2 WP-Based Prediction Algorithm

In the above algorithm, we can set C' as C because a newcoming MS can be queued until the resource is available after ASN GW relocation is completed. However, this approach cannot be applied to handover MSs because handover MSs are sensitive to handover latency. The acceptable handover delay is much less than the queuing delay of a new MS. Assuming that a handover MS arrives and C is reached. If the handover MS needs to wait for the ASN GW relocation of one Anchored MS, the handover latency will be too high.

We propose a prediction algorithm based on Wiener Process (WP) which provides a systematic way to determine when to request Anchored MSs to perform ASN GW relocation. In addition, the algorithm can also estimate how many Anchored MSs should be relocated.

Algorithm WP-based prediction algorithm

Require: At each t time interval.

1. If the number of samples is equal to k then
2. The oldest sample is discarded.
3. End if
4. $W(t)$ is recorded.
5. if $w(t) \geq T_{wnr}$ then
6. /* Stage 1: Generating the expected drift rate and the standard deviation rate */
7. $\hat{\mu}$ is computed by using k samples, $W(t-it)$, and (3).
8. $\hat{\delta}$ is computed by using k samples, $W(t-it)$, $\hat{\mu}$, and (4).
9. /* Stage 2: Estimating the number of MSs */
10. Computing ΔW by using $\hat{\mu}$, $\hat{\delta}$ and (2).
11. Computing $W(t+\Delta t)$ by using ΔW , $W(t)$, and (1).
12. /* Stage 3: Determining when and how many to perform ASN GW relocation */
13. If $\Delta W(t-\Delta t) > C$ then
14. $N \leftarrow \Delta W(t+\Delta t) + C + 1$
15. $N \leftarrow \min(n, N_a(t))$
16. Requesting n Anchored MSs to perform ASN GW relocation.
17. End if
18. End if

in the time interval $[t - i\tau - \tau, t - i\tau]$ are $W(t - i\tau) - W(t - i\tau - \tau)$, for $i = 0, \dots, k - 1$. Hence, μ can be estimated by $\hat{\mu}$

$$\begin{aligned} \hat{\mu} &= \frac{\sum_{i=0}^{k-1} (W(t - i\tau) - W(t - i\tau - \tau))}{k\tau} \\ &= \frac{W(t) - W(t - k\tau)}{k\tau}. \end{aligned} \quad (3)$$

Also, $\hat{\delta}$, the estimation of δ , is given by

$$\hat{\delta} = \sqrt{\frac{\sum_{i=0}^{k-1} (W(t - i\tau) - W(t - i\tau - \tau) - \hat{\mu}\tau)^2}{k\tau}}. \quad (4)$$

Wiener Process has been proven effective in modeling stochastic processes where the values of the random variables are affected by a large number of independent or weakly dependent factors, each with a relatively small impact. The $W(t)$ we want to model is impacted by a large number of factors. These factors are either independent or weakly dependent of each other. For example, $W(t)$ is impacted by the arrival rate of new MSs, arrival rate of handover MSs, average connection holding time, average network residence time, and so on. Based on the definitions and properties of Wiener Process, $W(t)$ is continuous and ΔW follows normal distribution.

3. Discussion

Comparing with traditional AC algorithms, the proposed GRAC decreases the blocking probability of new MSs in WiMAX networks. To see the reasons behind it, we use Flow (3) in Fig. 1 as an example. When deploying traditional AC algorithms without considering ASN GW relocation, a new incoming MS is blocked when $N_s(t) + N_a(t) = \min(T_{ncb}, C - N_h(t))$ in ASN GW A. However, there may be some Anchored MSs served in ASN GW A, that is, $N_a(t) > 0$. In contrast, in the proposed GRAC, ASN GW A will request one of the Anchored MSs to perform ASN GW relocation to relocate the traffic anchor point from ASN GW A to ASN GW B. Therefore $N_a(t)$ in ASN GW A is decreased. Thus, $N_s(t) + N_a(t) < \min(T_{ncb}, C - N_h(t))$.

Therefore, a new incoming MS can be accepted by ASN GW A. Besides, the $W(t)$ in ASN GW B is not increased.

4. Conclusion

In WiMAX standards, an ASN GW can decide when to perform ASN GW relocation. In this project, we consider that the system load is heavy, so Anchored MSs are forced to perform ASN GW relocation. We propose GRAC which considers admission control and ASN GW relocation jointly to improve the performance of WiMAX networks. The traditional AC algorithms cannot be used directly when the two-tiered mobility management is deployed in WiMAX because some MSs may be served by two ASN GWs. If there are many Anchored MSs, new incoming users will likely be rejected due to the lack of resources. In the proposed GRAC, the AC algorithm cooperates with the ASN GW relocation. When a new MS arrives and there is no resource for the newly arrived MS, the proposed GRAC will request an Anchored MS to perform ASN GW relocation. In addition, for handover MSs, the WP-based prediction algorithm can trigger the ASN GW relocation at an appropriate time.

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